# Géométrie fractale, algorithmes, calculabilité 

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## Iterated function systems (IFS)

Let $f_{1}, \ldots, f_{n}: \mathbb{R}^{d} \rightarrow \mathbb{R}^{d}$ be contracting maps.

## Theorem (Hutchinson 1981)

There is a unique nonemtpy compact set $X \subseteq \mathbb{R}^{d}$ such that

$$
X=f_{1}(X) \cup \cdots \cup f_{n}(X) .
$$

- Definition of "fractal"


## Example 1

- $f_{1}: x \mapsto \frac{1}{2} x$
- $f_{2}: x \mapsto \frac{1}{2} x+\binom{1 / 2}{0}$
- $f_{3}: x \mapsto \frac{1}{2} x+\binom{0}{1 / 2}$

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## Self-affine sets

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- Let $\mathcal{D} \subseteq \mathbb{Z}^{d}$ ("digits")


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Very particular kind of IFS:

- Affine maps given by integers
- Common matrix for all maps


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Meaning of "digits":

$$
X=\left\{\sum_{k=1}^{\infty} A^{-k} d_{k}:\left(d_{k}\right)_{k \geqslant 1} \in \mathcal{D}^{\mathbb{N}}\right\}
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## Self-affine tiles

## Definition

$\mathcal{D}$ is a standard set of digits if:

- $|\mathcal{D}|=|\operatorname{det}(A)|$
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## Theorem [Bandt, Lagarias-Wang]

Under these conditions:

- $X$ has nonempty interior
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- Many good properties and algorithms


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- Several other works in this direction...
- The real question: Is it decidable?


## Beyond self-affine tiles: affine IFS

- Several contractions $A_{1}, \ldots, A_{i}$ instead of just $A$
- Arbitrary affine mappings $f_{i}(x)=A_{i} x+v_{i}$


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## Beyond self-affine tiles: affine IFS

- Several contractions $A_{1}, \ldots, A_{k}$ instead of just $A$
- Arbitrary affine mappings $f_{i}(x)=A_{i} x+v_{i}$
- Several sets $X_{1}, \ldots, X_{n}$ instead of just $X$ (Graph-IFS)


## Graph-IFS (GIFS)

$$
\begin{aligned}
f_{1}(x) & =x / 2 \\
f_{2}(x) & =x / 2+\binom{1 / 2}{0} \\
f_{3}(x) & =x / 2+\binom{0}{1 / 2}
\end{aligned}
$$

$$
X=f_{1}(X) \cup f_{2}(X) \cup f_{3}(X)
$$



## Graph-IFS (GIFS)

$$
\begin{array}{ll}
f_{1}(x)=x / 2 & f_{4}(x)=x / 2+\binom{1 / 2}{1 / 2} \\
f_{2}(x)=x / 2+\binom{1 / 2}{0} & f_{5}(x)=x / 2+\binom{1 / 2}{3 / 4} \\
f_{3}(x)=x / 2+\binom{0}{1 / 2} & f_{6}(x)=x / 2+\binom{3 / 4}{3 / 4}
\end{array}
$$

$$
\left\{\begin{array}{l}
X=f_{1}(X) \cup f_{2}(X) \cup f_{3}(X) \cup f_{4}(Y) \\
Y=f_{5}(Y) \cup f_{6}(X)
\end{array}\right.
$$



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$$

$$
f_{f_{3}}^{f_{1}} C_{f_{6}}^{f_{4}}
$$



## Undecidability result

Theorem [J-Kari 2013]
For 2D affine GIFS with 3 states (with coefficients in $\mathbb{Q}$ ):

- $=[0,1]^{2}$ is undecidable
- empty interior is undecidable


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For 2D affine GIFS with 3 states (with coefficients in $\mathbb{Q}$ ):

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- Also true with diagonal $A_{i}$
- Computational tools: multi-tape automata


## Multi-tape automata

## $d$-tape automaton:

- alphabet $\mathcal{A}=A_{1} \times \cdots \times A_{d}$
- states $\mathcal{Q}$
- transitions $\mathcal{Q} \times\left(A_{1}^{+} \times \cdots \times A_{d}^{+}\right) \rightarrow \mathcal{Q}$


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\begin{aligned}
\mathcal{A} & =\{0,1\} \times\{0,1\} \\
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000
11

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0001

110

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> 000101
> 11000

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> 000101000
> 1100011

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$$

Accepted infinite word starting from $X$ :

> 0001010001
> 11000110

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$$

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> 0001010001000
> 1100011011

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\end{aligned}
$$

Accepted infinite word starting from $X$ :

$$
\begin{aligned}
& 0001010001000 \ldots \\
& 1100011011 \ldots
\end{aligned} \in \mathcal{A}^{\mathbb{N}}=(\{0,1\} \times\{0,1\})^{\mathbb{N}}
$$

## Multi-tape automaton $\longmapsto$ GIFS



## Multi-tape automaton $\longmapsto$ GIFS

$$
\begin{gathered}
u \mid v \\
\text { Transition } \\
\text { (tape alphabets } \left.A_{1}, A_{2}\right) \\
\text { Mapping } f\binom{x}{y}=\left(\begin{array}{cc}
\left|A_{1}\right|^{-|u|} & 0 \\
0 & \left|A_{2}\right|^{-|v|}
\end{array}\right)\binom{x}{y}+\binom{0 . u_{1} \ldots u_{|u|}}{0 . v_{1} \ldots v_{|v|}}
\end{gathered}
$$

## Multi-tape automaton $\longmapsto$ GIFS



Automaton:


## Multi-tape automaton $\longmapsto$ GIFS

(tape alphabets $A_{1}, A_{2}$ )

$$
\longmapsto \text { Mapping } f\binom{x}{y}=\left(\begin{array}{cc}
\left|A_{1}\right|^{|u|} & 0 \\
0 & \left|A_{2}\right|^{|v|} \mid
\end{array}\right)\binom{x}{y}+\binom{0 . u_{1} \ldots u_{|u|}}{0 . v_{1} \ldots v_{|v|}}
$$

## Automaton:



Associated GIFS:

$$
\left(\begin{array}{cc}
1 / 4 & 0 \\
0 & 1 / 4
\end{array}\right)\binom{x}{y}+\binom{0.01}{0.00} \text { ( }
$$

## Multi-tape automaton $\longmapsto$ GIFS

$$
\begin{aligned}
& \text { Transition } \overbrace{\bullet}^{u \mid v} \quad \begin{array}{c}
\text { (tape alphabets } A_{1}, A_{2} \text { ) } \\
\longmapsto \text { Mapping } f\binom{x}{y}=\left(\begin{array}{cc}
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\end{array}
\end{aligned}
$$

| Automaton | GIFS |
| :--- | :--- |
| states | fractal sets |
| edges | contracting mappings |
| \#tapes | dimension |
| alphabet $A_{i}$ | base- $\left\|A_{i}\right\|$ representation on tape $i$ |

## Multi-tape automaton $\longmapsto$ GIFS

$\rightarrow \overbrace{0}^{1011 \mid 11}$

- $f\binom{x}{y}=\left(\begin{array}{cc}1 / 16 & 0 \\ 0 & 1 / 4\end{array}\right)\binom{0 . x_{1} x_{2} \ldots}{0 . y_{1} y_{2} \ldots}+\binom{0.1011}{0.11}$


## Multi-tape automaton $\longmapsto$ GIFS

## 1011|11

$\rightarrow \xrightarrow{\infty}$

$$
\begin{aligned}
f\binom{x}{y} & =\left(\begin{array}{cc}
1 / 16 & 0 \\
0 & 1 / 4
\end{array}\right)\binom{0 . x_{1} x_{2} \ldots}{0 . y_{1} y_{2} \ldots}+\binom{0.1011}{0.11} \\
& =\binom{0.0000 x_{1} x_{2} \ldots}{0.00 y_{1} y_{2} \ldots}+\binom{0.1011}{0.11}
\end{aligned}
$$

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## 1011|11

$\rightarrow>$

$$
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- $\longrightarrow_{0}$

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& =\binom{0.1011 x_{1} x_{2} \ldots}{0.11 y_{1} y_{2} \ldots}
\end{aligned}
$$

## Key correspondence

GIFS fractal associated with automaton $\mathcal{M}$

$$
\left\{\binom{0 . x_{1} x_{2} \ldots}{0 . y_{1} y_{2} \ldots} \in \mathbb{R}^{2}:\binom{x_{1} x_{2} \ldots}{y_{1} y_{2} \ldots} \text { accepted by } \mathcal{M}\right\}
$$

## Multi-tape automaton $\longmapsto$ GIFS

## Example:



Automaton

$$
\frac{1}{2}\binom{x}{y}+\binom{0}{1 / 2}+\binom{0}{0}
$$



$$
=\left\{\binom{0 . x_{1} x_{2} \ldots}{0 . y_{1} y_{2} \ldots}:\left(x_{n}, y_{n}\right) \neq(1,1), \forall n \geqslant 1\right\}
$$

## Multi-tape automaton $\longmapsto$ GIFS

## Example:




## Multi-tape automaton $\longleftrightarrow$ GIFS

| Automaton | GIFS |
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| Languages properties | GIFS properties |
| :--- | :--- |
| $\exists$ configurations with $=$ tapes | Intersects the diagonal [Dube] |
| Is universal | Is equal to $[0,1]^{d}$ |
| Has universal prefixes | Has nonempty interior |
| $?$ | Is connected |
| $?$ | Is totally disconnected |
| Compute language entropy | Compute fractal dimension |

## Multi-tape automaton $\longleftrightarrow$ GIFS

## Theorem [Dube 1993]

It is undecidable if $X$ intersects the diagonal.

Proof idea:

- $X \cap\{(x, x): x \in[0,1]\} \neq \varnothing$
$\Longleftrightarrow$ Automaton accepts a word of the form $\binom{0 . x_{1} x_{2} \ldots}{0 . x_{1} x_{2} \ldots}$
- Reduce the Post-correspondence problem


## Language universality $\Longleftrightarrow$ nonempty interior

Fact 1: $\mathcal{M}$ is universal $\Longleftrightarrow X=[0,1]^{d}$

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Fact 2: $\mathcal{M}$ is prefix-universal $\Longleftrightarrow X$ has nonempty interior

- Example (universal with prefix 1 but not universal): one state, transitions 1, 10, 00 (one-dimensional)

$$
\begin{aligned}
& f_{1}(x)=x / 2+1 / 2 \\
& f_{2}(x)=x / 4 \\
& f_{3}(x)=x / 4+1 / 2
\end{aligned}
$$

$11 \square$


## Undecidability results

## Theorem [J-Kari 2013]

For 3-state, 2-tape automata:

- universality is undecidable
- prefix-universality is undecidable


## Corollary

For 2D affine GIFS with 3 states (with coefficients in $\mathbb{Q}$ ):

- $=[0,1]^{2}$ is undecidable
- empty interior is undecidable


## Conclusion \& perspectives

Decidability of nonempty interior:

|  | IFS | 2-state GIFS | $\geqslant 3$-state GIFS |
| :---: | :---: | :---: | :---: |
| dimension 1 | $?$ | $?$ | $?$ |
| dimension $\geqslant 2$ | $?$ | $?$ | Undecidable |

## Conclusion \& perspectives

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- dimension 1: we need new tools


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## References

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